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## Performance Evaluation and Configuration in Time-Sensitive Networking

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#### What is Time-Sensitive Networking?

- **Technically:** Collection of interacting IEEE standards
  - Still actively worked worked on
  - Few devices implement standards
- **Simplified**: Ethernet with QoS-guarantees
- Related Topics from our work
  - Autocorrelated traffic source modelling at Endpoints
  - Latency calculation on the Data Plane
  - Decentralized resource allocation on the Control Plane





### **Traffic source simulation**

- TSN networks modelled as queueing systems  $\rightarrow$  Autocorrelation important
- Our choice: Discrete Event Simulation
- ► We introduced **DARTA**<sup>[1]</sup>for modelling discrete stationary time-series with
  - Any discrete marginal distribution
  - Any autocorrelation structure
- Present for arrivals in
  - Industrial Networks
  - IoT traffic
  - Consumer applications
- Currently subject of research at our chair



### **DARTA - Functionality**

- ► Based on ARTA<sup>[2]</sup>mechanism
- Finds easily generatable Gaussian base process to transform to target process
- Main problem: Finding fitting autocorrelation for base process
- Main contribution is an integral approximation
- Mathematical proof is omitted here





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#### **DARTA - Results**

DARTA works!

- Distributions and Autocorrelations can be approximated with high accuracy
- Extreme cases can introduce difficulties
- Currently researching realistic usecases



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# LATENCY CALCULATION<sup>[3]</sup>

[3] Bounded Latency with Bridge-Local Stream Reservation and Strict Priority Queuing by Grigorjew, Alexej; Metzger, Florian; Hoßfeld, Tobias; Specht, Johannes; Götz, Franz-Josef; Chen, Feng; Schmitt, Jürgen in 11th International Conference on Network of the Future (2020)

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#### Where does Delay come from?



Model of a Switch



### Shaping

- Streams are shaped at each port
  - All streams destined for port arrive in queues
  - One queue for each of 8 priorities
  - FIFO queues
- Many shaping mechanisms available
  - Credit Based Shaping
  - Asynchronous Traffic Shaping
  - Strict Priority
  - ••••
- Consider strict priority for now
  - Selects highest priority available element
  - Does not pre-empt frames



Cable

#### **Delay Accumulation**

- Each stream *i* has a maximum burst rate  $r_i$  and a priority  $p_i$
- n-th device visited by *i* has
  - Maximum latency  $\delta_n^p$  per priority (configured)
  - Minimum processing time (hardware contraint)
- At each device/hop, a minimum/maximum accumulated delay can be computed
- We have some research on fitting choices for maximum latency





#### **Delay Accumulation**

- How much data from its own stream i can be in queue before packet A?
- Look at worst case:
  - Packet A has minimul delay (packets in front have less time processing)
  - Packet B is just leaving device 2 (device 2 has a latency bound)
  - Data in front could only be sent in between A and B

 $m_{device=n}^{stream=i} = \left(\operatorname{AccMaxD}_{n+1}^{i} - \operatorname{AccMinD}_{n}^{i}\right) \cdot r_{i}$ 

Also true for other streams of same priority



Time

#### **Delay Accumulation**

- How much data of higher priority stream h can be processed while packet A is in line?
- Same as before for when A arrives, but Packet B has latency guarantee  $\delta_n^{p_h}$
- But even after A is enqueued, higher priority traffic may arrive
- Only limited by latency guarantee δ<sup>p<sub>i</sub></sup><sub>n</sub> for A

$$m_{device=n, block=h}^{stream=i} =$$

$$\left(\operatorname{AccMaxD}_{n+1}^{h} - \operatorname{AccMinD}_{n}^{h} + \delta_{n}^{p_{i}}\right) \cdot r_{i}$$

#### Device 2 Egress Queue (High Priority)









#### **Summary of Delay Computation**

► What is the total delay for a stream I at device n?

▶ Need to consider all streams of higher priority  $H_{p_i}$  or equal priority  $E_{p_i}$ 

$$\delta_{device=n}^{stream=i} = \sum_{h \in H_{p_i}} \frac{m_{device=n,block=h}^{stream=i}}{r_n} + \sum_{e \in E_{p_i}} \frac{m_{device=n}^{stream=e}}{r_n} + \left( \max_{l \in L_{p_i}} \frac{l}{r_n} \right)^2$$

- > Packet from lower priority streams  $L_{p_i}$  may be transmitted when packet A arrives
- Formula only relies on previously computed values
- Can be used for decentralized delay computation





## **DECENTRALIZED RESOURCE ALLOCATION**





#### **Classification of TSN Models**

Centralized vs Decentralized

- Centralized system has single controller doing all work
- Decentralized system requires nodes to
  - Allocate resources
  - Maintain routing tables
  - Propagate information
- But it is *probably* more resilient

#### Static vs Dynamic

- Static: System can be configured optimally before use
- Dynamic: Traffic sources and QoS requirements are updated on the fly

But it *seems* more flexible

- **Question:** How to allocate resources in decentralized, dynamic system?
- We currently work on simulation of the Resource Allocation Protocol<sup>[4]</sup>

[4] Draft Standard for Local and Metropolitan Area Networks — Bridges and Bridged Networks — Amendment: Resource Allocation Protocol

#### **Resource Allocation Protocol Simulation**

- Simple publish-subscribe model
- Device 0 sends talker announce (TA)
- Bridges check QoS constraints
- Gets broadcast through the system
- Device 2 subscribes with listener attach (LA)
- Devices check constraints and allocate resources
- When LA arrives at device 0, configuration is completed successfully



#### **Check failures**

Delay bound checks can occur in 2 places

- During TA propagation
  - Device computes delay check failure
  - Flag is set on TA
  - TA is broadcast further
- During LA propagation
  - Device computes delay check failure
  - Flag is set on LA
  - LA device is informed of failure
  - LA is sent further
- Process is not entirely standardized



#### Conclusion

- ► We covered three aspects of our TSN research
- Traffic source modelling
  - Success highly dependent on distribution and autocorrelation
  - We are currently testing for realistic parameter sets
- Delay computation
  - More complex shapers available
  - Some require time synchronization → Comes with own challenges (and protocols)
- Decentralized dynamic TSN
  - Approach is greedy
  - Still largely untested in the wild
  - May only work well with certain topologies
  - Standards still in development

#### Sources

[1] DARTA: Generation of Autocorrelated Random Numbers using Discrete AutoRegression To Anything by Geißler, Stefan; Raunecker, David; Lange, Stanislav; Hossfeld, Tobias at *ITC 35th -Networked Systems and Services (2023)* 

[2] Autoregressive to anything: Time-series input processes for simulation by Cario, Marne C.; Nelson, Barry L. in *Operations Research Letters (1996)* 

**[3] Bounded Latency with Bridge-Local Stream Reservation and Strict Priority Queuing** by Grigorjew, Alexej; Metzger, Florian; Hoßfeld, Tobias; Specht, Johannes; Götz, Franz-Josef; Chen, Feng; Schmitt, Jürgen in *11th International Conference on Network of the Future (2020)* 

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